TOPOLOGICAL ANALYSIS OF SUBSTITUTION TILING SPACES

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QUESTION.

Suppose you had an **INFINITE SUPPLY** of blue and white SQUARE TILES. How might you form an infinite tiling of the plane \mathbb{R}^2 ?

AN INTERESTING ANSWER.

Use a tiling substitution!



AN EXAMPLE OF A *substitution rule*.



Iterate the substitution to get arbitrarily large patches:



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- ▶ In this talk we'll see that substitution tiling spaces are Cantor set fiber bundles that can be seen as inverse limits and that their cohomology can be computed.

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- Standing assumptions: finite local complexity, nonperiodicity

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- σ can be applied to any patch of tiles by applying σ to each tile t in the patch and placing the result atop L(t).
- ▶ We call $\sigma(t)$ a supertile, $\sigma^2(t)$ a 2-supertile, and $\sigma^n(t)$ an *n*-supertile

HALF-HEX SUBSTITUTION RULE

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The half-hex substitution rule.

The prototile set \mathcal{P} contains six tile types rather than one because our framework considers tiles the same only if they are translates of one another.

A FEW HALF-HEX SUPERTILES 2-, 3-, and 4-supertiles



 $\sigma^4(A)$

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The set of all n-supertiles acts as the 'language' of Ω
If T ∈ Ω, then T − x ∈ Ω for any translation x ∈ ℝ^d

The big ball metric

How to measure the distance between tilings

Let T and T' be tilings of \mathbb{R}^d from a prototile set \mathcal{P} .

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DEFINITION. Let R(T,T') be the supremum of all $r \ge 0$ such that there exists $\vec{x}, \vec{y} \in \mathbb{R}^d$ with

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We define

$$d(T,T') = \min\left\{\frac{1}{R(T,T')},1\right\}$$

SUBSTITUTION TILING SPACES: TOPOLOGICAL BASICS

LEMMA. If Ω is of finite local complexity, then Ω is complete and compact.

LEMMA. Under mild conditions, Ω is connected. Each tiling in Ω defines a path component that is homeomorphic to \mathbb{R}^d , and there are uncountably many path components.

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- Simplicial, singular, and cellular cohomology don't work well either, since they study path connected components.
- Čech cohomology is computable and we will do the half-hex.

Cylinder sets

To visualize any $T' \in B_{\epsilon}(T) \subset \Omega$, take $T \cap B_{1/\epsilon}(0)$, a big central patch in T.

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THEOREM (SW)

A tiling space that satisfies certain $WLOG^1$ hypotheses is a fiber bundle over the torus, with totally disconnected fiber.

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 - Since every (n + 1)-supertile is composed of *n*-supertiles, the CW complex for the (n + 1)-supertiles maps onto that of the *n*-supertiles.
- Every tiling $T \in \Omega$ corresponds to an element of the inverse limit by noting the location of 0 in each of its *n*-supertiles.

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We make Γ_n in exactly the same fashion, using *n*-supertiles. This gives instructions for tiling larger and larger regions.

(SUBSTITUTION) TILING SPACES AS INVERSE LIMITS Forcing the border

A substitution forces the border if there is an N such that every N-supertile determines the tiles immediately adjacent to it.



The half-hex substitution forces the border with N = 2.

(SUBSTITUTION) TILING SPACES AS INVERSE LIMITS The 'forgetful' map $\phi_n : \Gamma_{n+1} \to \Gamma_n$



THE FORGETFUL MAP φ₁ ON PART OF Γ₂
Each (n + 1)-supertile is composed of n-supertiles, and φ_n: Γ_{n+1} → Γ_n is a continuous cellular map.

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(SUBSTITUTION) TILING SPACES AS INVERSE LIMITS The inverse limit formalism

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- ► Elements of $\lim_{\leftarrow} (\Gamma_n, \phi_n)$ give instructions for making tilings:
 - ▶ p_0 tells what tile to place at the origin, and precisely where
 - \triangleright p_1 tells what supertile to place around that tile
 - \triangleright p_2 tells what 2-supertile to place around the 1-supertile, etc.

THEOREM (AP)

When the substitution forces the border, Ω and $\lim_{\leftarrow} (\Gamma_n, \phi_n)$ are homeomorphic.

(If it doesn't force the border we use a trick called "collaring")

(WE NEED IT BUT CAN AVOID COMPUTING IT DIRECTLY)

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 - $\check{H}^*(\Gamma_n, \mathbb{Z}) = H^*(\Gamma_n, \mathbb{Z})$ because each Γ_n is a CW complex

What we do in practice

To compute the Čech cohomology of the tiling space Ω of a substitution σ :

• If σ does not force the border, use a collaring trick to ensure that Ω and $\lim_{\leftarrow}(\Gamma_n, \phi_n)$ are homeomorphic.

What we do in practice

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- Figure out how φ^{*} acts on H^{*}(Γ, Z). For each dimension 0, 1, ..., d the result is a matrix.
- Take the direct limit of the matrix to get the cohomology of the inverse limit and thus of Ω.

TOP ČECH COHOMOLOGY OF THE HALF-HEX

COHOMOLOGY OF APPROXIMANTS



The labelled CW complex Γ

After a bit of linear algebra we obtain that

$$H^0(\Gamma, \mathbb{Z}) = \mathbb{Z}$$
 $H^1(\Gamma, \mathbb{Z}) = \mathbb{Z}^2$ $H^2(\Gamma, \mathbb{Z}) = \mathbb{Z}^3$

Let $A^*, B^*, C^*, D^*, E^*, F^*$ represent the dual cochains to the 2-chains A, B, C, D, E, F. The equivalence relation in the quotient for H^2 gives $A^* = D^*, B^* = E^*$, and $C^* = F^*$.

Generators for $H^2(\Gamma, \mathbb{Z})$ are A^*, B^*, C^*

TOP ČECH COHOMOLOGY OF THE HALF-HEX Forgetful map as substitution

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The forgetful map ϕ on 1-chains is computed to be:

 $\begin{array}{lll} A \rightarrow A + C + D + E & B \rightarrow B + D + E + F & C \rightarrow A + C + E + F \\ D \rightarrow A + B + D + F & E \rightarrow A + B + C + E & F \rightarrow B + C + D + F \end{array}$

TOP ČECH COHOMOLOGY OF THE HALF-HEX

PULLBACK OF THE FORGETFUL MAP

We compute the pullback $\phi^* : C^2(\Gamma_n) \to C^2(\Gamma_{n+1})$:

$$\phi^*(\omega)(q) = \omega(\phi(q)), \ \omega \in C^2(\Gamma_n) \text{ and } q \in C_2(\Gamma_{n+1})$$

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Consider $A^* \in C^2(\Gamma_n)$ and any $q \in C_2(\Gamma_{n+1})$. $\blacktriangleright \phi^*(A^*(q)) = A^*(\phi(q)) = 0$ if $\phi(q)$ contains no A, \Longrightarrow

q = B and F

$$\begin{array}{ccc} A \xrightarrow{\phi} A + C + D + E & B \xrightarrow{\phi} B + D + E + F & C \xrightarrow{\phi} A + C + E + F \\ D \xrightarrow{\phi} A + B + D + F & E \xrightarrow{\phi} A + B + C + E & F \xrightarrow{\phi} B + C + D + F \end{array}$$

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- For all other choices of q, $A^*(\phi(q)) = 1$ since $\phi(q)$ contains one copy of A

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- For all other choices of q, $A^*(\phi(q)) = 1$ since $\phi(q)$ contains one copy of A
- Thus $\phi^*(A^*) = A^* + C^* + D^* + E^* \cong 2A^* + B^* + C^*$

$$A \xrightarrow{\phi} A + C + D + E \quad B \xrightarrow{\phi} B + D + E + F \quad C \xrightarrow{\phi} A + C + E + F$$
$$D \xrightarrow{\phi} A + B + D + F \quad E \xrightarrow{\phi} A + B + C + E \quad F \xrightarrow{\phi} B + C + D + F$$

TOP ČECH COHOMOLOGY OF THE HALF-HEX The direct limit of ϕ^*

Using similar logic for the other two generators of $H^2(\Gamma_n, \mathbb{Z})$ we find that in the basis A^*, B^*, C^*, ϕ^* acts as the matrix $\begin{pmatrix} 2 & 1 & 1 \\ 1 & 2 & 1 \\ 1 & 1 & 2 \end{pmatrix}$ $\check{H}^2(\Omega, \mathbb{Z}) = \lim_{\longrightarrow} (H^2(\Gamma, \mathbb{Z}), \phi^*) = \lim_{\longrightarrow} \left(\mathbb{Z}^3, \begin{pmatrix} 2 & 1 & 1 \\ 1 & 2 & 1 \\ 1 & 1 & 2 \end{pmatrix} \right)$

The eigenvalues are 4, 1, 1 and in the final analysis we arrive at:

$$\check{H}^2(\Omega,\mathbb{Z}) = \mathbb{Z}[1/4] \oplus \mathbb{Z} \oplus \mathbb{Z}$$

Top Čech cohomology of the half-hex

INTUITIVE INTERPRETATION

 $\check{H}^2(\Omega,\mathbb{Z}) = \mathbb{Z}[1/4] \oplus \mathbb{Z} \oplus \mathbb{Z}$

 $1 \in \mathbb{Z}[1/4]$ is the cochain that knows when it sees a tile $1/4 \in \mathbb{Z}[1/4]$ is the cochain that knows when it sees a supertile $1/4^n \in \mathbb{Z}[1/4]$ is the cochain that knows when it sees an *n*-supertile

The two copies of \mathbb{Z} are generated by cochains that can tell an A from a B and a B from a C.

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The two copies of \mathbb{Z} are generated by cochains that can tell an A from a B and a B from a C.

Taken together the cohomology has the ability to recognize types of supertiles of all orders, up to the identification $A^* = D^*, B^* = E^*$, and $C^* = F^*$.

In a similar fashion we can compute that

$$\check{H}^1(\Omega,\mathbb{Z}) = \lim_{\longrightarrow} (H^1(\Gamma,\mathbb{Z}),\phi^*) = \lim_{\longrightarrow} \left(\mathbb{Z}^2, \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix} \right)$$

And ultimately

$$\check{H}^1(\Omega,\mathbb{Z}) = \mathbb{Z}[1/2] \oplus \mathbb{Z}[1/2].$$

This is harder to interpret but reflects the linear expansion factor of two on edges.

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